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(11) **EP 1 251 641 A2** 

(12)

### **EUROPEAN PATENT APPLICATION**

(43) Date of publication:

23.10.2002 Bulletin 2002/43

(51) Int CI.7: H03M 5/14

(21) Application number: 02252790.7

(22) Date of filing: 19.04.2002

(84) Designated Contracting States:

AT BE CH CY DE DK ES FI FR GB GR IE IT LI LU

MC NL PT SE TR

Designated Extension States:

AL LT LV MK RO SI

(30) Priority: 20.04.2001 KR 2001021360

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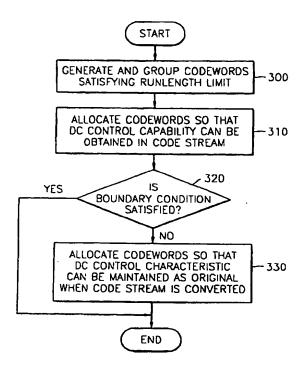
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### (54) Code generation and allocation method

(57)A method for generating and allocating codewords is provided. The method includes allocating one of two selectable codewords b1 and b2 as codeword b when a preceding codeword a and a following codeword b form a code stream X, in which codewords b1 and b2 have opposite INVs which are parameters indicating whether the number of '1s' contained in a codeword is an odd number or an even number and when the code stream of a and b1 is X1, and the code stream of a and b2 is X2, allocating codewords such that the INVs of X1 and X2 are maintained to be opposite when a or b1(b2) should be replaced by another codewords in compliance with a predetermined boundary condition given between codewords. According to the method, by using a short codeword having less bits as a main conversion codeword, high efficiency is achieved in recording density. Also, when codewords which do not satisfy the run length conditions are replaced by other codewords, the codewords are allocated so that the DC suppression capability of the code stream can be maintained, and therefore higher DC suppression capability of the code stream is provided.

FIG. 3



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### Description

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[0001] The present invention relates to generation and allocation of modulation codes of source codes to be recorded on a recording medium, and more particularly, to a method for generating and allocating codewords in which codewords having a restricted run length are generated and the generated codewords are allocated so that the DC control characteristic of a code stream is maintained.

[0002] In a Run Length Limited (RLL) code represented by (d, k, m, n), the performance of a code is evaluated mainly based on the recording density and the capability to suppress the DC component of the code. Here, m denotes the number of data bits (the number of so-called source data bits, which is also referred to as the number of information word bits), n denotes the number of codeword bits after modulation (the number of so-called channel bits), d denotes the minimum number of a series of '0s' that can exist between '1' and '1' in a codeword, and k denotes the maximum number of a series of '0s' that can exist between '1' and '1' in a codeword. The interval between bits in a codeword is represented by T.

[0003] In a modulation method, to improve recording density it is used to reduce the number of codeword bits n while regarding d and m as given conditions. In the RLL code, however, d which is the minimum number of a series of '0s' that can exist between '1' and '1' in a codeword, and k which is the maximum number of a series of '0s' that can exist between '1' and '1' in a codeword, should be satisfied. If, with this (d, k) condition satisfied, the number of data bits is m, the number of codewords satisfying RLL(d, k) should be equal to or greater than 2<sup>m</sup>. Moreover, in order to actually use this code, run length constraints, that is, RLL(d, k) conditions, should be satisfied in a part where a codeword is linked to another codeword. In addition, when the DC component of a code affects the system performance, it is desirable to use a code which has a DC suppression capability.

[0004] The main reason for suppressing the DC component in the RLL modulated code stream is to minimize a reproducing signal's affect on a servo band. Hereinafter, methods for suppressing the DC component will be referred to as Digital Sum Value (DSV) control methods.

[0005] DSV control methods can be broadly classified into two types. One is a method having a DSV control code itself, where the DSV control code is capable of controlling a DSV. The other is a method of inserting a merge bit at each DSV control time. An Eight to Fourteen Modulation plus (EFM+) code performs DSV control using a separate code table, while an EFM code or a (1, 7) code performs DSV control by inserting a merge bit.

[0006] Therefore, the shape of the prior art modulation code group having the DSV control code itself capable of controlling suppression of the DC component and satisfying the conditions described above is as shown in Figure 1, in which each of a predetermined number of main conversion code groups has a corresponding code group for controlling suppression of the DC component. Each main conversion code group and its corresponding code group form a pair so that the DC component can be suppressed and controlled. In this case, there are some characteristics that distinguish codewords of the predetermined main conversion code groups. That is, there are no identical codewords between the main conversion code groups A and B. If duplicated codes are used, there might be the conversion code groups C and D for demodulating the duplicated codes, where there are no identical codewords between the conversion code groups C and D, but codewords in the code group A or B may be in the code group C or D for demodulating duplicated codes. The number of codewords in the main conversion code groups A and B and the conversion code groups C and D for demodulating duplicated codes is 2<sup>m</sup> if the number of bits in the source word before conversion is m. [0007] If code groups E through H are DC suppression control code groups used for suppressing DC components together with code groups A through D, respectively, the characteristics of codewords in each of the code groups E through H are the same as the characteristics of codewords in the main code groups A through D respectively. That is, the same conditions for generating duplicated codewords or the same conditions for determining the number of lead zeros in a codeword are applied to each of the DC suppression control code groups E through H for controlling suppression of DC components and the conversion code groups A through D.

[0008] For example, the characteristics of the EFM+ code, which is used in current Digital Versatile Discs (DVD), has a run length condition of RLL(2, 10) and a codeword length (n) of 16 bits, is as shown in Figure 2. The main conversion code groups are MCG1 ("A" in Figure 1) and MCG2 ("B" in Figure 1) and the conversion code groups for demodulating duplicated codes are DCG1 ("C" in Figure 1) and DCG2 ("D" in Figure 2). There are four DSV code groups ("E~H" in Figure 1) which make pairs with respective conversion code groups to control suppression of DC components. There are no identical codewords between the four conversion code groups and the four DSV code groups which are code groups for controlling DC components.

[0009] Also, the conditions for generating duplicated codewords in the entire code groups are the same, and the characteristics of codewords in each code group pair that can control DC components (MCG1 and the first DSV code group, MCG2 and the second DSV code group, DCG1 and the third DSV code group, or DCG2 and the fourth DSV code group) are the same.

[0010] That is, a codeword having a continuous sequence of from 2 to 5 zeros from the Least Significant Bit (LSB) of the codeword is generated using duplicated codewords. This rule is applied to each code group in the same manner.

In each of the codewords of the first DSV code group for controlling suppression of DC components, which controls suppression of DC components together with the main conversion code group MCG1, there is a continuous sequence of between 2 and 9 '0s' from the Most Significant Bit (MSB). In each of the codewords of the second DSV code group for controlling suppression of DC components, which controls suppression of DC components together with the main conversion code group MCG2, there is either 0 or 1 '0' continuing from the MSB. Some bits (here, b15(MSB) or b3) in the codewords of the third DSV code group for controlling suppression of DC components, which controls suppression of DC components together with the conversion code group DCG1 for demodulating duplicated codes are '0b', while some bits (here, b15(MSB) or b3) in the codewords of the fourth DSV code group for controlling suppression of DC components, which controls suppression of DC components together with the conversion code group DCG2 for demodulating duplicated codes, some bits (here, b15(MSB) and b3) are '1b". In developing 8 to 15 modulation code which has an advantage in the recording density aspect compared to the prior art modulation method EFM+ which uses the modulation code group shown in Figure 1 or 2, the original characteristics of a code stream change when a change occurs in a codeword because of a boundary rule applied to the locations adjacent to a boundary which connects a codeword to another codeword.

[0011] It is an aim of the present invention to provide a method for generating and allocating codewords in which a codeword having a run length restriction is generated and the codeword is allocated so that the original characteristics of a code stream are maintained without change even when a codeword is replaced according to the boundary rule when a code stream is allocated.

[0012] According to a first aspect of the present invention there is provided a method for generating and allocating codewords of source words which are to be recorded on a recording medium, the method including generating codewords satisfying predetermined run length conditions and grouping codewords according to each run length condition; and allocating the codewords such that a code(word) for the source word is capable of controlling suppression of DC components.

[0013] It is preferable that when a predetermined boundary condition is not satisfied in the code stream, allocating codewords such that codewords which satisfy the boundary condition and maintain the DC control characteristics which are considered when the initial codewords are allocated replace the initial codewords.

[0014] It is preferable that the step for generating codewords includes generating codewords satisfying the length of a predetermined first codeword, and predetermined run length conditions, grouping the codewords according to each predetermined run length condition to generate a main conversion codeword table; generating DC control codewords satisfying the length of a predetermined second codeword, and predetermined run length conditions in order to control DC components in the code(word) stream, grouping the DC control codewords, and to generate a code conversion table for controlling DC components; and generating additional DC control codewords by taking codewords which satisfy the predetermined run length conditions and are not needed in the main conversion codeword table, and grouping the additional DC control codewords.

[0015] According to a second aspect of the present invention there is provided an allocation method for allocating codewords generated for source words to be recording on a recording medium, the method including when a preceding codeword a and a following codeword b form a code stream X, allocating one of two selectable codewords b1 and b2 as codeword b, in which codewords b1 and b2 have opposite INVs which are parameters indicating whether the number of '1s' contained in a codeword is an odd number or an even number and when the code stream of a and b1 is X1, and the code stream of a and b2 is X2, allocating codewords such that the INVs of X1 and X2 are maintained to be opposite when a or b1(b2) should be replaced by another codewords in compliance with a predetermined boundary condition given between codewords.

[0016] It is preferable that when the predetermined boundary condition is that the number of continuous '0s' between codewords should be at least 2, and when the number of continuous '0s' from the Least Significant Bit (LSB) of the codeword a in the Most Significant Bit (MSB) direction is 0, and the number of continuous '0s' from the MSB of the codewords b1 or b2 in the LSB direction is 1, code changes of either the codeword a or b1(b2) occur to satisfy the boundary condition.

[0017] It is preferable that when the number of continuous '0s' between the codewords a and b is 1 or 0, the codeword a or b is changed such that the number of 0s forming the boundary is greater than 2 and less than 10.

[0018] It is preferable that the codeword a of the code stream X1 and the codeword a of the code stream X2 are changed to other codewords such that the resulting codewords a of code streams X1 and X2 have the same INV value, and as a result, by the INVs of codewords b1 and b2 following the codewords a respectively, the INVs of the X1 and X2 become different.

[0019] According to a third aspect of the present invention there is provided an allocation method for allocating codewords of source words to be recording on a recording medium, the method including when a preceding codeword b and a following codeword c form a code stream Y, allocating one of two selectable codewords b1 and b2 as the codeword b, wherein codewords b1 and b2 have opposite INVs which are parameters indicating whether the number of '1s' contained in a codeword is an odd number or an even number and when the code stream of b1 and c is Y1, and

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the code stream of b2 and c is Y2, allocating codewords such that INVs of Y1 and Y2 are maintained to be opposite when the codeword b1, b2 or c should be replaced by another codeword in compliance with a predetermined boundary condition between codewords.

[0020] It is preferable that when the predetermined boundary condition is that the number of continuous '0s' between codewords should be at least 2, and when the number of continuous '0s from the Least Significant Bit (LSB) of the codeword c toward the Most Significant Bit (MSB) is 1, codeword b which does not satisfy the boundary condition and is xxxxxxxxxxxx1001 or xxxxxxxxxxx10001 appears in both b1 an b2.

[0021] For a better understanding of the invention, and to show how embodiments of the same may be carried into effect, reference will now be made, by way of example, to the accompanying diagrammatic drawings in which:

Figure 1 is a diagram of an example of the shape of a prior art modulation code group;

Figure 2 is a table showing the prior art code group and the characteristics of codewords included in the code group;

Figure 3 is a flowchart showing a method for generating and allocating codes according to the present invention;

Figure 4 is a table showing a variety of codeword groups of main conversion tables and the characteristics of codewords in each codeword group;

Figure 5 is a table showing a variety of codeword groups of a conversion table for DC control and the characteristics of codewords in each codeword group;

Figure 6 is a table showing a variety of codeword groups of an auxiliary conversion table for DC control and the characteristics of codewords in each codeword group;

Figure 7 is a diagram for showing what should be considered for the run length conditions when codewords a and b are connected;

Figure 8 is a table showing an example of changes in INV before code conversion and after code conversion when the run length conditions described in Figure 7 are not satisfied;

Figure 9 is a diagram showing an example of code stream branching due to selective codewords b1 and b2 for DC control;

Figure 10 is a graph showing the relationship between INV values of a code stream pair;

Figures 11a through 11e are main conversion code according to the present invention;

Figures 12a through 12j are code conversion tables for DC control according to the present invention;

Figures 13a and 13b are auxiliary code conversion tables for DC control according to the present invention; and

[0022] Figure 14 is a graph showing the difference between the frequency spectrum when codewords of the code conversion table for DC control according to the present invention are used in 25% of all of the codewords, and the frequency spectrum when prior art EFM+ modulation codewords are used.

[0023] Figure 3 is a flowchart showing a method for generating and allocating codes according to a preferred embodiment of the present invention. According to the method for generating and allocating codewords of source words to be recorded on a recording medium, codewords satisfying predetermined run length conditions are generated and the generated codewords are grouped according to each run length condition in step 300. The codewords are allocated so that the code(word) streams for source words are capable of controlling DC components in step 310. It is determined whether or not predetermined boundary conditions are satisfied in the code stream in step 320. If the conditions are not satisfied, the codewords are replaced by codewords satisfying the boundary conditions while the DC control characteristics which are considered when the original codewords are allocated can be kept.

[0024] Code tables of the codewords for source code conversion are roughly divided into three types: 1) main conversion tables, 2) conversion tables for controlling DC components, and 3) auxiliary conversion tables for controlling DC components.

[0025] Figure 4 is a table showing a variety of codeword groups of main conversion tables and the characteristics of codewords in each code group. It is assumed that d denotes the minimum run length limit of a codeword, k denotes

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the maximum run length limit of a codeword, m denotes the number of bits of source data, n denotes the number of bits of a codeword after modulation, End Zero (EZ) denotes the number of '0s' in a continuous sequence from the LSB of a codeword in a direction toward the MSB of the codeword, and LZ denotes the number of '0s' in a continuous sequence from the MSB of a codeword in a direction toward the LSB of the codeword. For example, codewords that satisfy d=0, k=10, m=8, n=15,  $0 \le EZ \le 8$  are divided according to the following LZ conditions:

- 1) number of codewords satisfying 2≤LZ≤10: 177
- 2) number of codewords satisfying 1≤LZ≤9: 257
- 3) number of codewords satisfying 0≤LZ≤6: 360
- 4) number of codewords satisfying 0≤LZ≤2: 262

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[0026] If the number of bits of source data satisfies m=8, the number of codewords for conversion should be 256 or more. However, in condition 1), the number of codewords does not amount to 256. Therefore, the number of codewords in condition 1) can amount to 256 by taking some codewords from a condition having surplus number of codewords. In this case, 83 codewords from the codewords satisfying group 3)'s LZ condition may be taken and added to group 1). Then, the numbers of codewords included in conditions 1) through 4) are 260, 257, 277(=360-83), and 262, respectively, and satisfy the minimum number of modulation codewords, that is, 256 for 8-bit source data. In the table of Figure 4, Main Code Group 1 (MCG1) is the name of a code group containing codewords satisfying condition 1) and some (83) codewords are taken from codewords satisfying condition 3). MCG2 and MCG4 are the names of codewords satisfying condition 2), and 4), respectively. MCG3 is the name of codewords satisfying condition 3), excluding the 83 codewords taken by MCG1. In each of the main code groups MCG1 through MCG4, only 256 codewords can be used as conversion codes for source codes.

[0027] Figure 5 is a table showing a variety of codeword groups of a conversion table for DC control and the characteristics of codewords in each codeword group. For example, assuming that d=2, k=10, m=8, n=17, and 0≤EZ≤8, conversion code tables for controlling DC components may include the following 4 groups (corresponding to DCG1, DCG2, DCG3, and DCG4 of Figure 5, respectively) according to the LZ conditions:

- 1) number of codewords satisfying 2≤LZ≤10: 375
- number of codewords satisfying 1≤LZ≤9: 546
- 1) number of codewords satisfying 0≤LZ≤6: 763
- 1) number of codewords satisfying 0<LZ<2: 556

[0028] Each group forming a conversion table for controlling DC components should have at least 2 codewords that selectively correspond to one source data, and therefore should have at least  $512 \ (= 2^8 + 2^8)$  codewords for 8-bit source data. Since the number of codewords in the code group satisfying the LZ condition 1) is less than 512, code group 1) can take surplus codewords from other code groups satisfying other LZ conditions to amount to the number of 512. For example, in the above embodiment, code group 1) may take 177 codewords from the code group satisfying the condition 3) so as to have  $552 \ (=375 + 177)$  codewords.

[0029] Figure 6 is a table showing a variety of codeword groups of an auxiliary conversion table for DC control and the characteristics of codewords in each code group. For example, among codewords satisfying d=2, k=10, m=8, and n=15, codewords satisfying 9≤EZ≤10, the remaining codewords of the main code conversion groups (MCGs), and codewords satisfying LZ=7, 8 or LZ=4, 5 are used as codewords of auxiliary code groups (ACGs) for controlling suppression of DC components. The conditions for generating these codewords will now be explained in detail. The following conditions correspond to ACG1 through ACG4, respectively, which are names of the auxiliary conversion tables for controlling suppression of DC components:

- 1) 5 codewords (satisfying 9≤EZ≤10 and LZ≠0) + the remaining 4 codewords (in the MCG1) = 9 codewords,
- 2) 5 codewords (satisfying 9≤EZ≤10 and LZ≠0) + 1 remaining codewords (in the MCG1) = 6 codewords,
- 3) 5 codewords (satisfying  $9 \le EZ \le 10$  and  $LZ \ne 1$ ) + 15 codewords (satisfying  $7 \le LZ \le 8$  and  $0 \le EZ \le 8$ ) = 41 codewords, the remaining 4 codewords in the MCG1 = 9 codewords,
- 4) 7 codewords (satisfying 9≤EZ≤10 + the remaining 6 codewords in the MCG4) + 85 codewords (satisfying 3≤LZ≤5 and 0≤EZ≤8) = 98 codewords.

[0030] When codeword a and codeword b are connected, the junction where the two codewords are connected should satisfy a run length (d, k) condition. Figure 7 is a diagram showing what should be considered for the run length conditions when codewords a and b are connected. Satisfying the run length condition means that in Figure 7 a value obtained by adding the end zero (EZ\_a) of codeword a and the lead zero (LZ\_b) of codeword b is equal to or greater than the minimum run length d and equal to or less than the maximum run length k.

[0031] Figure 8 is a table showing an example of changes in INV (whose meaning will be described below) before code conversion and after code conversion when the run length conditions described in Figure 7 are not satisfied. Codeword b is determined in a group indicated by the EZ of the preceding codeword, codeword a. When a or b is included in a code group which does not have enough codewords to meet the condition and takes codewords from other code groups, the (d, k) condition may not be satisfied. In this example, the EZ of codeword a changes to satisfy the run length condition, which is referred to as the boundary rule. Variable INV which indicates whether the number of '1s' in a codeword stream is an even number or an odd number may change from the previous INV while the codeword a didn't change (unclear), according to the boundary rule. Due to this characteristic, attention should be paid to allocation of a codeword between code conversion tables capable of controlling suppression of DC components.

[0032] Figure 9 is a diagram showing an example of code stream branching due to selective codewords b1 and b2 for DC control. One of the major features of the code conversion of the present invention is that the codewords of two code conversion tables that can be selected for DC control are allocated so that they have opposite INV characteristics. When the previous INV changes according to the boundary rule as described above, if the INVs of both codewords in the two code conversion tables that can be selected change, then there will be no problem. Otherwise, characteristics of codewords having opposite INV cannot be maintained. For this reason, a code conversion table is designed considering the following.

[0033] First, in A of Figure 9, that is, at the junction where the codeword a and the codeword b are connected to each other, if b1 and b2, which can be selected as codeword b, are codewords in DCG11 and DCG12, respectively, which are regrouped in the code conversion table DCG1 shown in Figure 5 after separating codewords which correspond to the same source code but have different INVs, or if b1 and b2 are codewords of MCG1 and MCG2, respectively, then codewords in which LZ\_b1 (the number of LZs of codeword b1) and LZ\_b2 (the number of LZs of codeword b2) is 1 are allocated on the location. By doing so, when the EZ of the codeword a is '0', according to the boundary rule, the INV of codeword a changes in both the code stream containing the codeword b1 and the code stream containing the codeword b2, or the INV of codeword a does not change in either the code stream containing the codeword b1 or the code stream containing the codeword b2, such that the INVs of the two code streams are maintained to be opposite. An example is as follows:

27 224 250 source data stream1(before conversion)000001000010001(MCG3) 30 code 000001000001001 (MCG1) 010010010000000 (MCG1) conversion) 000001000010001 stream1(after code 35 010010010000000 000001000001000 0 1 INV1 stream2 (before conversion) 000001000010001 (MCG3) code 40 000001000001001 (MCG1) 01001000000000 (ACG1) conversion) 000001000010001 stream2(after code 0100100000000000 000001000001000 45 1 1 1 INV2

[0034] Next, in B of Figure 9, that is, at the junction where codeword b and codeword c are connected to each other, if codewords b1 and b2 are respectively included in code conversion tables DCG11 and DCG12, DCG21 and DCG22, DCG31 and DCG32, DCG41 and DCG42, MCG1 and ACG1, MCG2 and ACG2, MCG3 and ACG3, or MCG4 and ACG4, and (xx)xxxxxxxxxx1001 or (xx)xxxxxxxxxx10001, INV may change according to the boundary rule due to the LZ of the following codeword c. Therefore, these codewords bland b2 are allocated to the location for corresponding same source data in each table such that the INVs of the two code streams are maintained to be opposite. An example is as follows:

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	source	e data	250	. 15:	2	210	
5	code s	streaml(before	conve	ersion)	000	001000010001 (MCG	3)
·	010000	000010001001 (DC	CG11)	00000010	0000001	(MCG1)	
	code	stream1(after	conv	version)		00000100001000	00
10	010000	000010001001		0000010	0000001		
	INV1		0		0	0	
	code s	stream2(before	conve	ersion)	000	001000010001 (MCG	3)
15	010010	000010001001 (DC	G12)	01000000	1001001	(MCG1)	
	codes	tream2(after	conv	ersion)		000001000010000	00
	010010	000010001001		01000000	1001001		
20	INV2		0		1	1	

[0035] For the junctions A and B of Figure 9, the codewords are first allocated to the location corresponding same source data in each code conversion table (DCG11 and DCG12 or MCG1 and ACG1) considering above. Referring to the following example, in point B, according to the boundary rule, the INVs of code stream1 and code stream2 are maintained to be opposite and the INVs of code stream3 and code stream4 are maintained to be opposite and the INVs of code stream1 and code stream3 are maintained to be opposite and the INVs of code stream4 are maintained to be opposite.

source data 250 152 7 code stream1(before conversion) 000001000010001 (MCG3) 35 0100000010001001(DCG11) 010000010010001(MCG1) code stream1(after conversion) 000001000010000 0100000010001000 010000010010001 40 INV1 0 1 1

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	code	stream2(before	conve	ersion)	00000	01000010001 (MCG3)
5	01000	000010001001 (DC	G11)	01001001001	L0001 (A	ACG1)
	code	stream2(after	conv	ersion)		0000010000100000
	01000	000010001000		01001001001	10001	
10	INV2		0		1	0
	code	stream3 (before	conve	ersion)	00000	01000010001 (MCG3)
	01001	.000010001001 (DC	G12)	01000001001	L0001 (M	(CG1)
15	code	stream3(after	conv	ersion)		000001000010000
	01001	.000010001000		01000001001	10001	
	EV/I		0		0	0
20	code	stream4 (before	conve	ersion)	00000	01000010001 (MCG3)
	01001	.000010001001 (DC	G12)	01001001003	10001 (2	ACG1)
	code	stream4(after	conv	ersion)		0000010000100000
25	01001	.000010001000		0100100100	10001	
	INV4		0		0	1

[0036] As described above, considering changes in the INV of a codeword due to the boundary rule in a codeword stream, codewords are allocated so that the INV polarities of a codeword pair after modulation is always be maintained to be opposite. Figure 10 is a graph showing the relationship of INV values of this code stream pair. If codewords are allocated such that the INV values of a code stream pair are always opposite, a codeword can be selected so that a code stream which is DC components between the code stream pair is formed.

[0037] Exceptions to the rule that INV values are maintained to be opposite at point A of Figure 9 may occur when source data is from 251 to 255 in the code conversion table for controlling DC components. In such exceptional cases, the CSV signs of codewords are made to be opposite so that the difference between DSV values in the code stream pair is made.

[0038] Figures 11a through 11e are main conversion code tables in which codewords are generated and allocated considering conditions described above.

[0039] Figures 12a through 12j are code conversion tables for DC control in which codewords are generated and allocated considering conditions described above.

[0040] Figures 13a and 13b are auxiliary code conversion tables for DC control in which codewords are generated and allocated considering conditions described above.

[0041] Figure 14 is a graph showing the difference between the frequency spectrum when codewords of the code conversion table for DC control according to the present invention are used in 25% of all of the codewords, and the frequency spectrum when prior art EFM+ modulation codewords are used. The graph shows that in a low frequency band, the frequency spectrum of the modulated code stream according to the present invention is almost the same as the frequency spectrum of the EFM+, which indicates that the code stream of the present invention has almost the same capability of suppressing DC components as that of the EFM+ method.

[0042] Accordingly, since the present invention uses 15-bit codes as the main conversion code and selectively uses 17-bit DC control codes for controlling DC components, the present invention has better efficiency in recording density than the prior art EFM+ code and has the same DC suppression capability as the EFM+ code.

[0043] In the present invention, by using a short codeword having less bits as a main conversion codeword, high efficiency is achieved in recording density.

[0044] Also, when codewords which do not satisfy the run length conditions are replaced by other codewords, the codewords are allocated so that the DC suppression capability of the code stream can be maintained, and therefore higher DC suppression capability of the code stream is provided.

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[0045] The reader's attention is directed to all papers and documents which are filed concurrently with or previous to this specification in connection with this application and which are open to public inspection with this specification, and the contents of all such papers and documents are incorporated herein by reference.

[0046] All of the features disclosed in this specification (including any accompanying claims, abstract and drawings), and/or all of the steps of any method or process so disclosed, may be combined in any combination, except combinations where at least some of such features and/or steps are mutually exclusive.

[0047] Each feature disclosed in this specification (including any accompanying claims, abstract and drawings), may be replaced by alternative features serving the same, equivalent or similar purpose, unless expressly stated otherwise. Thus, unless expressly stated otherwise, each feature disclosed is one example only of a generic series of equivalent or similar features.

[0048] The invention is not restricted to the details of the foregoing embodiment(s). The invention extends to any novel one, or any novel combination, of the features disclosed in this specification (including any accompanying claims, abstract and drawings), or to any novel one, or any novel combination, of the steps of any method or process so disclosed.

### Claims

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 A method for generating and allocating codewords of source words which are to be recorded on a recording medium, the method comprising:

generating codewords satisfying predetermined run length conditions and grouping codewords according to each run length condition; and

allocating the codewords such that a codeword for the source word is capable of controlling suppression of DC components.

2. The method of claim 1, further comprising:

when a predetermined boundary condition is not satisfied in the code stream, allocating codewords such that codewords which satisfy the boundary condition and maintain the DC control characteristics which are considered when the initial codewords are allocated replace the initial codewords.

- 3. The method of claim 1 or 2, wherein in order to make code streams capable of controlling suppression of DC components allocating codewords such that a pair of codeword streams having opposite INV characteristics are made selectable, where INV indicates whether the number of '1s' is an odd number or an even number.
- 4. The method of claim 1, 2 or 3, wherein the step for generating codewords comprises:

generating codewords satisfying the length of a predetermined first codeword, and predetermined run length conditions, grouping the codewords according to each predetermined run length condition to generate a main conversion codeword table;

generating DC control codewords satisfying the length of a predetermined second codeword, and predetermined run length conditions in order to control DC components in the code(word) stream, grouping the DC control codewords, and to generate a code conversion table for controlling DC components; and

generating additional DC control codewords by taking codewords which satisfy the predetermined run length conditions and are not needed in the main conversion codeword table, and grouping the additional DC control codewords.

- 5. The method of claim 4, wherein when the bit length of the source word is 8, the length of a codeword in the main conversion code table is 15 bits.
- 6. The method of claim 5, wherein the main conversion code table contains groups of codewords, the groups formed of a group of codewords each having from 2 to 10 Lead Zeros (LZs), a group of codewords each having from 1 to 9 LZs, a group of codewords each having from 1 to 6 LZs, and a group of codewords each having from 0 to 2 LZs, while having from 0 to 8 End Zeros (EZs).

- 7. The method of claim 6, wherein among the groups of the main conversion code table, a group having a lesser number of codewords than the minimum number of codewords for converting the source data takes surplus codewords from a group having a greater number of codewords than the minimum number of codewords so as to amount to the minimum number of codewords.
- 8. The method of claim 5, wherein the length of the codewords of the DC control conversion code table is 17 bits.
- 9. The method of claim 8, wherein the DC control conversion code table contains groups of codewords, the groups formed of a group of codewords of which LZ is from 2 and to 10, a group of codewords of which LZ is from 1 to 9, a group of codewords of which LZ is from 0 to 6, and a group of codewords of which LZ is from 0 to 2, while EZ is from 0 to 8.
- 10. The method of claim 9, wherein each code group of the DC control conversion table has as much codewords as a source word can correspond to pairs of codewords, each pair of codewords has opposite INV characteristics and is selectable.
- 11. The method of claim 10, wherein among the groups of the DC control conversion code table, a group having a lesser number of codewords than the minimum number of needed codewords takes surplus codewords from a group having a greater number of codewords than the minimum number of codewords, so as to have equal to or greater than the minimum number of codewords.
- 12. The method of claim 8, wherein the auxiliary DC control conversion table having a group of codewords, each having a length of 15 bits, having from 9 to 10 EZs, and having at least one LZ, and codewords which are taken from surplus codewords of the first main conversion code group;
  - a group of codewords, each having a length of 15 bits, having from 9 to 10 EZs, and having at least one LZ, and codewords which are taken from surplus codewords of the second main conversion code group;
  - a group of codewords, each having a length of 15 bits, having from 9 to 10 EZs, and having no LZ, the surplus codewords of the third main conversion code group, and having 7 to 8 LZs, or having 0 to 8 EZs;
  - a group of codewords, each having a length of 15 bits, and having 9 or 10 EZs, the surplus codewords of the fourth main conversion code group, and having 3 to 8 LZs and 0 to 8 EZs.
- 13. The method of claim 12, wherein when the code stream pairs are a, b1, c and a, b2, c, respectively, and b1 and b2 are DC control codewords having opposite INV characteristics, codewords are allocated such that the INV characteristics of code streams after conversion are maintained to be opposite even if code changes of a, b1, b2, or c occur due to violation of the predetermined run length between a, and b1(b2) or b1(b2) and c.
- 40 14. An allocation method for allocating codewords generated for source words to be recording on a recording medium, the method comprising:
  - when a preceding codeword a and a following codeword b form a code stream X, allocating one of two selectable codewords b1 and b2 as codeword b, wherein codewords b1 and b2 have opposite INVs which are parameters indicating whether the number of '1s' contained in a codeword is an odd number or an even number; and
  - when the code stream of a and b1 is X1, and the code stream of a and b2 is X2, allocating codewords such that the INVs of X1 and X2 are maintained to be opposite when a or b1(b2) should be replaced by another codewords in compliance with a predetermined boundary condition given between codewords.
  - 15. The method of claim 14, wherein when the predetermined boundary condition is that the number of continuous '0s' between codewords should be at least 2, and when the number of continuous '0s' from the Least Significant Bit (LSB) of the codeword a in the Most Significant Bit (MSB) direction is 0, and the number of continuous '0s' from the MSB of the codewords b1 or b2 in the LSB direction is 1, code changes of either the codeword a or b1(b2) occur to satisfy the boundary condition.
  - 16. The method of claim 14, wherein when the number of continuous '0s' between the codewords a and b is 1 or 0,

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the codeword a or b is changed such that the number of 0s forming the boundary is greater than 2 and less than 10.

- 17. The method of claim 16, wherein the codeword a of the code stream X1 and the codeword a of the code stream X2 are changed to other codewords such that the resulting codewords a of code streams X1 and X2 have the same INV value, and as a result, by the INVs of codewords b1 and b2 following the codewords a respectively, the INVs of the X1 and X2 become different.
- **18.** An allocation method for allocating codewords of source words to be recording on a recording medium, the method comprising:

when a preceding codeword b and a following codeword c form a code stream Y, allocating one of two selectable codewords b1 and b2 as the codeword b, wherein codewords b1 and b2 have opposite INVs which are parameters indicating whether the number of '1s' contained in a codeword is an odd number or an even number; and

when the code stream of b1 and c is Y1, and the code stream of b2 and c is Y2, allocating codewords such that INVs of Y1 and Y2 are maintained to be opposite when the codeword b1, b2 or c should be replaced by another codeword in compliance with a predetermined boundary condition between codewords.

- 19. The method of claim 18, wherein when the predetermined boundary condition is that the number of continuous '0s' between codewords should be at least 2, and when the number of continuous '0s' from the Least Significant Bit (LSB) of the codeword c toward the Most Significant Bit (MSB) is 1, codeword b which does not satisfy the boundary condition and is xxxxxxxxxxxx1001 or xxxxxxxxxxx10001 appears in both b1 and b2.
- 25 20. The method of claim 18, wherein when the number of continuous '0s' between the codewords a and b is 1 or 0, the codeword a or b is changed such that the number 0s forming the boundary is greater than 2 and less than 10.

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FIG. 1

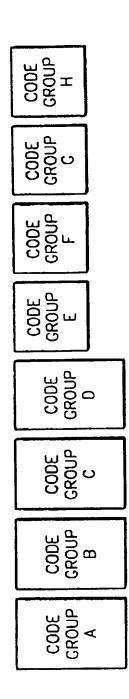


FIG. 2

NCC				2		3		4
CODE GROUP CONTAINING CODEVORD INDICATED BY	CONVERSION CODE GROUP	OC SUPPRESSION CODE GROUP FOR OC SUPPRESSION CONTROL WITH THE HELP OF MCCI	CONVERSION CODE GROUP	OC SUPPRESSION CODE GROUP FOR OC SUPPRESSION CONTROL WITH THE HELP OF MCG2	CONVERSION CODE GROUP	DC SUPPRESSION CODE GROUP FOR DC SUPPRESSION CONTROL WITH THE HELP OF DCC1	CONVERSION CODE GROUP	CONVERSION CODE GROUP FOR CONVERSION CODE GROUP FOR
אנפ	MCG1	1st DSV CODE GROUP	MCC2	2nd DSV CODE GROUP	1920	3rd DSV CODE GROUP	7900	4th DSV CODE CROUP
CHARACTERISTIC   LZ=2~9	6~2=7	6~2=71	17=0~1	1~0=Z∩	b15 (MSB)=b3=0	)=63=0	b15 (MSB):	b15 (MSB)= 1 EE b3=1
METHOD FOR DUPUCATED CODE		W3ODO	ROD OF EZ=2	~5 REPEATEDLY OC	CURS IN ALL M	CODEWROD OF EZ=2~5 REPEATEDLY OCCURS IN ALL KINDS OF CODE GROUPS	UPS	

FIG. 3

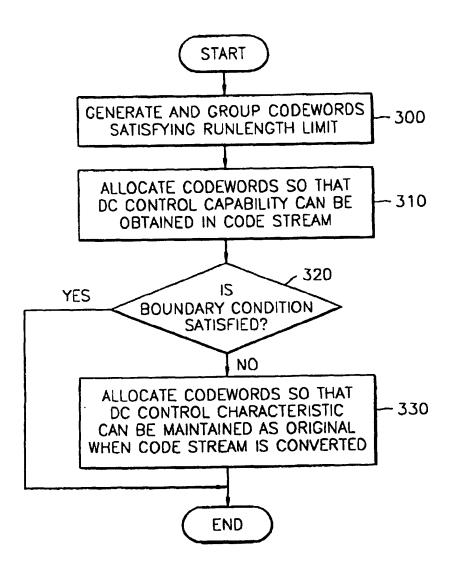


FIG. 4

	LZ (EZ)	NUMBER	ADD	NUMBER	DELETE	NUMBER	TOTAL NUMBER (DUPLICATED)
MCG1	2≤LZ ≤10 (0≤EZ≤8)	177	LZ=1 (0≤EZ≤8)	83			260(4)
MCG2	1 <u>≤</u> LZ <u>≤</u> 9 (0≤EZ <u>≤</u> 8)	257			·	·	257(1)
MCG3	0≤LZ≤6 (0≤EZ≤8)	360			LZ=1 (0≤EZ≤8)	83	277(21)
MCG4	0≤LZ≤2 (0≤EZ≤8)	262					262(6)

FIG. 5

\(\frac{71}{2}\)	LZ . (EZ)	NUMBER	ADD	NUMBER	DELETE	NUMBER	TOTAL NUMBER
DCG1	2≤LZ≤10 (0≤EZ≤8)	375	LZ=1 (0≤EZ≤8)	177			552
DCG2	1≤LZ≤9 (0≤EZ≤8)	546					546
DCG3	0≤LZ≤6 (0≤EZ≤8)	763			LZ=1 (0≤EZ≤8)	177	586
DCG4	0≤LZ≤2 (0≤EZ≤8)	556					556

FIG. 6

	LZ (EZ)	NUMBER	ADD	NUMBER		NUMBER	TOTAL NUMBER
ACG1	LZ≠0 (9≤EZ≤10)	5	SURPLUS CODE OF MCG1	4			9
ACG2	LZ≠0 (9≤EZ≤10)	5	SURPLUS CODE OF MCG2	1			6
ACG3	LZ≠1 (9≤EZ≤10)	5	SURPLUS CODE OF MCG3	21	7≤LZ≤8 (0≤EZ≤8)	15	41
ACG4	(9≤EZ≤10)	7	SURPLUS CODE OF MCG4	6	3≤LZ≤5 (0≤EZ≤8)	85	98

FIG. 7

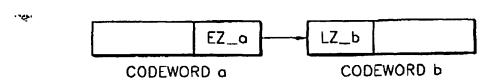


FIG. 8

CODEWORD a	CODEWORD 6	CHANGE IN INV
xxxxxxxxxx001001 (BEFORE CONVERSION) xxxxxxxxxx001000 (AFTER CONVERSION)		
xxxxxxxxx010001 (BEFORE CONVERSION) xxxxxxxxx010000 (AFTER CONVERSION)		CHANGE
xxxxxxxxx100001 (BEFORE CONVERSION) xxxxxxxxx100100 (AFTER CONVERSION)	0100xxxxxxxxxx	
~		NO
xxx10000000001 (BEFORE CONVERSION) xxx100000000100 (AFTER CONVERSION)		

FIG. 9

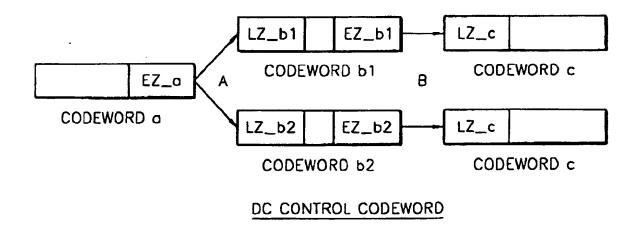
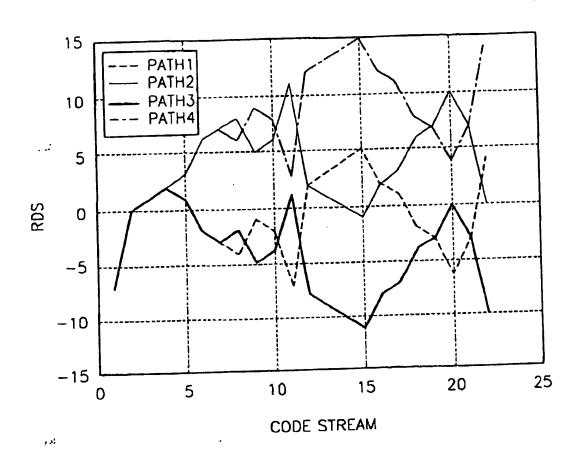


FIG. 10



## FIG. 11A

DATA	NCG1		NCC:		MCG3		NCG4	<del></del>
DATA SYMBOL	Code Word MSB LSB	NCG	Code Word MSB LSB	NCG	Code Word MSB LSB	NCG	Code Word MSB LSB	NCG
000	000100100000000	4	0001001000000000	4	100100000000001	1	100100000000001	1
000	000100100000000	4	000010010000000	1	001001000000001	j	0100100000000001	li
002	0010001001000000	4	000100100100000	4	0000001000000001	l i l	001001000000001	li
002	010001001000000	4	000000100000000	4	100000000100001	i	1001001000000001	l i
003	010000100100000	4	000000010000000	4	100000001000001	i	100010010000001	l i
005	010000001001001	li	000000000100010	2	1001000000000010	2	100001001000001	. 1
006	010000010010000	3	0001000000000001	1	0010010000000010	2	100000100100001	1
007	010000010010001	lī	0100100100000001	i	0010000000000010	2	100100010000001	1
008	010010000010010	2	010001001000001	1	100100000010010	2	100010001000001	1
009	010010010000001	1	010000100100001	1	001001000010010	2	100001000100001	1
010	000001001000000	4	0000100000000001	1	000100100010010	2 2	0100100100000001	1
011	000000100100000	4	0100100010000001	1	100000000010010	2	0100010010000001	1
012	000000010010000	3	010001000100001	1	100000000100010	2	010000100100001	1
013	000000001001000	3	001001001000001	1 1	1001000000000100	3	100000000100001	1
014	000000000100100	3	001000100100001	1	001001000000100	3	100100000000010	2
015	000000000100010	2	0000010000000001	1 ! !	000100100000100	3	010010000000010	1 2
016	000000000100001	1	010010000100001	1 1	000000100010000	3	001000000000010	1 4
017	000000000010000	3	001001000100001	1	100000010010000	3	100100100000010	1 2
018	010001001000001	1	000100100100001	)	100000001001000	3	100010010000010	1 5
019	010010001000000	4	0000001000000001	1	100000000100100	3	100001001000010	1 5
020 021	010010000000010	2	0010010000000001	1   1	10000010001000	3	10000010010010	1 2
021	010000100100001	4	00010010000000001	l i l	100000001000100	3	100100010000010	l ž
023	010010000000001	li	0000100100000001	l i l	100001000010000	Š	100010001000010	<u>2</u>
024	0001000000000001	l î	000001001000001	l i l	100000100001000	3	100001000100010	2
025	001001001000000	1 4	000000100100001	i	10000000100000	4	100000100010010	2
026	000000000010001	li	0100010000000001	1	000100100000000	4	010010010000010	2
027	010010010000000	4	001000100000001	1	000010010000000	4	010001001000010	2
028	0000100000000001	1	0001000100000001	1	000001001000000	4	100100000010010	22222222222222222223
029	010000001001000	3	000010001000001	l l	000000100100000	4	010010000010010	2
030	010010001000001	1	000001000100001	1 1	100010010000000	4	100000000010010	1 5
031	010000000100100	3	010000100000001	!	100001001000000	4	100000000100010	1 5
032	010001000100001	1 3	001000010000001	1 1	100000100100000	4 4	100000001000010	1 5
033 034	010000100010000	2	000010000100001	l i l	100100010000001001	i	010000000100010	l ž
034	010000000001001	í	010000010000001	l i l	0010010000001001	l i	100100000000100	3
<b>036</b> :	001001000000001	î	001000001000001	l i l	000100100001001	l i	100000000001000	3 -
037	001000100000000	4	000100000100001	lī	100000010010001	1 1	010000000000100	3
038	001000000000010	2	010000000100001	1	100000001001001	1	100000000010000	3
039 🚮	~0001001000000001	1	010000001000001	1	100000100010001	1	010000000001000	3 3 3
040	000100010000000	4	001000000100001	1	100000010001001	1	001000000000100	3
041	010000100010001	1	000000000100001	1	100010010000001	1	100100100000100 100010010000100	١١
042	000010010000001	1	000000001000001	1	100001001000001	1	100001001000100	3 3 3 3
043	000000100010000	3	000000010000001	1 2	100000100100001	1 1	10000100100100	1 3
044 045	010010000100000 0000010010000001	4	0010000000000010	2	1000100001000001	li	010000000010000	1 3
045	0000000100100001	3	0100100100000010	2	100001000100001	l i l	001000000001000	3
047	000000100100001	1	010001001000010	2	000010000000001	l i l	100100100001000	3
048	000000001000100	3	010000100100010	2	100100001000001	l i l	100100010000100	3
049	000000001000010	2	010000010010010	2	100010000100001	l i l	100010010001000	3
050	010010000001001	ī	000010000000010	2	001001001000001	1 1	100010001000100	3
051	010010000000100	3	010010001000010	2	001000100100001	1	100100010010000	3
052	000010001000000	4	010001000100010	2	0000010000000001	1	100100001001000	3
053	0001000000000010	2	010000100010010	2	100100000100001	1	100100000100100	3
054	000001000100000	4	001001001000010	2	001001000100001	1 1	010010010010000	3
055	010000010001001	1	001000100100010	2	000100100100001	1	100000010010000	ا ا
056	000001000010000	3	001000010010010	2	000100100000001	1	100000001001000	3
057	001001000100000	4	000001000000010	2	000010010000001	] ]	100000000100100	3
058	000000100001000	3	010010000000010	2	000001001000001	1	100000100010000	3
059	000000010000100	3	001001000000010	2			10000001000100	3
060	0000000000010010	2	0100000010	2	100010000000001	1	100000001000100	

## FIG. 11B

SARBOL   SCHOOL   S		NCG1		NCG2		MCC3		NCG4	
1000000000000000000000000000000000000		Code Word	NUC.		NCG		NCG		NCG
	31 MOUL		_			1100	-, -		3
		001001000001001	-			00100010000001			3
0.000   0.00		00100100100000100				000010001000001	1 1		
066  001001001000001   2		001001001000001			2	000001000100001			1 -
066		000100100100000	4		2			0.10000.100000000	1 :
Control   Cont		0010010000000010	2		2	00100001000000		100100101000000	1 -
		010001001000010		000100010000010	2			10070010010000000	
Control   Cont				00001000100010			- 1		4
					2	001000001000001			1
					2	000100000100001	1		
0.0000010001001001   2			1	001000010000010	2	100000010000001	, -	100001000100000	1 :
074   01000001001001   2   0000010001001   2   10000010010001   2   10000001000100   4   076   010000001001001   3   01000000100101   2   10000010010010   2   1001000000100000   4   077   00001000000001   1   00100000010010   2   01000000010010   2   100000001000000   4   078   0000010000000   1   001000000100010   2   0010000000000			3		2	001000000100001		100001000100000	
076								0100010010000000	
078		0010000001001001			5	10000100100010	2	010000100100000	
079						100000010010010	2		
0010000100100010		000001000000001		001000000010010	2				
B86				010000001000010			2		
OBS   ODIO0000001000   1   ODIO000000100   2   ODIO0000000000   1   ODIO000000000   1   ODIO000000000   1   ODIO000000000   1   ODIO000000000   2   ODIO000000000   1   ODIO000000000   2   ODIO000000000   1   ODIO000000000   2   ODIO000000000   2   ODIO000000000   1   ODIO000000000   2   ODIO0000000000   2   ODIO000000000   2   ODIO000000000   2   ODIO0000000000   2   ODIO0000000000   2   ODIO000000000   2   ODIO0000000000   2   ODIO000000000   2   ODIO0000000000   2   ODIO0000000000   2   ODIO0000000000   2   ODIO0000000000   2   ODIO0000000000   2   ODIO00000000000000000000000000000000000							151		
OBS   ODIO0000001000   1   ODIO000000100   2   ODIO0000000000   1   ODIO000000000   1   ODIO000000000   1   ODIO000000000   1   ODIO000000000   2   ODIO000000000   1   ODIO000000000   2   ODIO000000000   1   ODIO000000000   2   ODIO000000000   2   ODIO000000000   1   ODIO000000000   2   ODIO0000000000   2   ODIO000000000   2   ODIO000000000   2   ODIO0000000000   2   ODIO0000000000   2   ODIO000000000   2   ODIO0000000000   2   ODIO000000000   2   ODIO0000000000   2   ODIO0000000000   2   ODIO0000000000   2   ODIO0000000000   2   ODIO0000000000   2   ODIO00000000000000000000000000000000000						00000100010010	2		1
			13		2		2		
0.000000000000000			1 3		2	100010000100010	2		
0.00000010000000			1		2		2	01000000000001001	
088		001000000100100			2		5	00100000000010001	
088 001000000001001					2		1 2	100100100001001	1
090   01000000001000   3   00001010101010   2   00010010000010   2   00100000010010   1   1   1   1   1		001000000010010				000001000000010	2	100010010001001	
092   0100100000001   1   010010000100   2   0000010010010   2   10000010010001   1   1000000100100   2   10000010010010   1   10000010010010   2   10000010010010   1   100000100100010   2   10000010010001   1   100000100100010   2   100000100100010   1   100000100100010		010000000001001		000010010010010	2	000100100000010	2		, -
093			2		2	000010010000010	1 2		
093	092				2	000000100100000	2	100000001001001	1 -
095   00000100000010   2 00100100010010   2 0001001000010   2 01000001001001   1 000100100010010   1 00010001			3	010001000010010	1 5	100010000000000000000000000000000000000	2	100000100010001	
096   00100100000001   1   0001010010010   2   0001001000010   2   0100000100001				001001000100010	2	001000100000010	2	100000010001001	
100		0010010010000001			2	000100010000010			
100		010010000100100	3	000100010010010	2		1 2	0100000001001001	li
100	098	010010000100001			1 2		1 5	10010000100001	
101   0010010010000		7 000000100100010			1 3		1 2	010010001000001	
102			3	001000000000000000000000000000000000000	3		2	010001000100001	
103		001001001000000		01000000010000	13		) Z	001001001000001	
104		001001000010010		001000000001000	] 3		2	1001000100100001	
105   00010010001000   3	104			0001000000000100	1 3		2	010010000100001	1
107   000010001000001   1   010001001001000   3   100010010010010   2   010011000000001   1   1000   01000100			3	010001001000100		001000000010010	2	001001000100001	
108				010000100100100	3	100010010010010		1000100000000001	
109				001000000010000	3	100000010000010	2	0100010000000001	
110		001000010000000				001000000100010	1 2		1 .
111		010000001000010		000010000000100	1 3	000100000010010	1 5	0100001000000001	
113		010000000100001		0100100000000100	13	001000010000010	1 2	001000010000001	1
113   00100100100100   3   0000100100100   3   000010000010010   2   01000001000001   1   115   01000100100100   2   00000100100100   3   000010010010010   2   001000001000001   1   16   00000100100100   1   000000100100100   3   10010000010010   2   010000001000001   1   17   00010010010010   3   01001000001000   3   10010000010010   2   010000001000001   1   18   010010000100	112	0010000000000100			3	00100000100010	2	100000100000001	
115	113	1 000001000100100100	1 4	000010010000100	3	000010000010010	2	010000010000001	
116			2	000001001000100	3	000010010010010			
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		000000100010001	1	000000100100100	13				
118   010001000010001   1   010001000001000   3   00100010001001   2   010000001000001   1   119   0000000100001	117	000100100100100		010010000001000	13	10001000010010		100000010000001	
119   0000001000101   1   001001000100   2   001100100101   2   001100000100001				01000100000100	3	001000100010010	j   2	010000001000001	
190 I WINNINGSTREET   ( ) [IIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIII	119 120	000000010001001	3	001000100000100	1 3			001000000100001	11

## FIG. 11C

DATA	MCG1		JICG2		NCG3		MCG4	
SYMBOL	Code Word	NCG	Code Word	NCG	Code Word	NCG	Code Word	NCG
-	MSB LSB	↓	I CON	-	MSB LSB		NSB LSB	
121 122	000000000100000	4 2	000100100001000	3	000100010010010	2	010000100100010	2
123	000100100010010	ĺi	000010010001000	3	000000100000010 100000000001000	2	010000010010010 100100001000010	2 2
124	000010010000100	i	000010001000100	3	10000000001000	3	10010000100010	5
125	010000010000100	3	000001001001000	3	001000000000100	1 3	100001000010010	2
126	000100001000000	4	000001000100100	3	100010010000100	3	010010001000010	l ž
127	001000000100010	2	01000001001(000	3	100001001000100	,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,	010001000100010	2
128	001000000010001	1	010000001001000	3	100000100100100	3	010000100010010	2
129 130	010000000010000	3	010000000100100	3	001000000001000	3	001001001000010	2
131	010000100001000	2	010000010001000	3	000100000000100 100010010001000	3	001000100100010 001000010010010	1 5
132	000000100000100	3	010000001000100	3	1000100010001000	3	0010010000000010	1 5
133	010000100001001	Ĭ	001000010010000	3	100001001001000	3	100010000000010	1 5
134	010010010001000	3	001000001001000	3	100001000100100	3	0100010000000010	l ž
135	000000001000000	4	001000000100100	3 3	001000000010000		001000100000010	2
136	000010010010010	2	010001000010000	3	000100000001000	3	100001000000010	2
137 138	000001001001001	3	010000100001000	3 3	000010000000100	3	010000100000010	1 2
139	000100000010010	2	001000100010000	3	000010010000100 000001001000100	ოოოო	001000010000010	5
140	0001000000001001	۱ĩ	001000010001000	l 3 l	00000100100100	131	100010010010010	2
141	001000000001000	3	001000001000100	3	100100000001000	3	100000010000010	Ž
142	0010010010000010	2	000100010010000	3	100010000000100	I 3 I	010000001000010	2
143	001001000100001	1	000100001001000	3	001001000001000	3	001000000100010	2
144 145	001000100010000 000000010010010	3 2	000100000100100 010010000010000	3	001000100000100	3	100100010010010	222222222222222222222222222222222222222
146	010001001001001	ĩ	010001000001000	3	000100100001000 000100010000100	3	010010010010010	5
147	0001000000000100	Э	010000100000100	3	00001001000100	3	010000010000010	ĺž
148	001000100100010	2	001001000010000	3	000010001000100	3	001000001000010	2
149	001000100010001	1	001000100001000	3	000001001001000	3	001001000010010	2
150 151	001000010001000	3	001000010000100	3	000001000100100	3	100100000100010	2
151	0100001000000001	1 3	000100100010000 000100010001000	3	100000010000100	3	100010000010010	2
153	001000010010010	2	000100001000100	3	001000001001000	3	010001000010010	2
154	001000010001001	1	000010010010000	3	001000000100100	3	001001000100010	2
155	001000001000100	3	000010001001000	3	100010000010000	33333	001000100010010	2
156	000010000100000	4	000010000100100	3	100001000001000	3	100001001001000	3
157 158	010000010000010 010000001000001	2	000000010010000	3	100000100000100	3	100001000100100	3
159	010000000100001	4	000000001001000	3	001000100010000	3	010010010000100	3
160	000100010010000	3	000000100010000	3	001000001000100	∣ัลั∣	010000100100100	3
161	000000100100100	3	000000010001000	3 3	000100010010000	3 3 3	001000000010000	3
162	000100100100001	1	000000001000100	3	000100001001000	3	010010000000100	3
163	000100001001000	3	000001000010000	3	000100000100100	3	0010010000000100	3
164 165	010001000000010 001000001000010	2 2	000000100001000	3	100100000010000	3	100100000001000	3
166	001000001000010	î	010010010010000	3	10001000000100	3	010010000001000	3
167	001000000100001	î l	010010001001000	ă	001001000010000	3	010001000000100	ž
168	010010001000100	3	010010000100100	3 3	001000100001000	3	001001000001000	3
169	010001001000100	3	001001001001000	3	001000010000100	3	001000100000100	3
170	000000001000001	1	001001000100100	3	000100100010000	3	100001000010000	3
171	000100000100100	3	000100100100100	3	000100010001000	3	100000100001000	3
172 173	010010000001000	3 3	000010000010000	3	000100001000100	3 3	100000010000100	3
174	0010001000000010	2	000001000001000	9	000010010010000	3	010000010001000	93333
175	000100000100010	2	010010010001000	ğ	000010000100100	3	010000001000100	3
176	000100001000001	ĭ	010010001000100	3	000001000010000	3	001000010010000	3
177	000100000010001	1	010001001001000	3	000000100001000	3	001000001001000	3
178	010001001001000	3	010001000100100	3	100100001001000	3	001000000100100	3
179	010000100100100	3	001001001000100	3	100100000100100	9	100010000010000	3
180	010001000000100	-3-1	001000100100100	<u> </u>	001001001001000	_3_1	TOTOTOTOTOTOTO	<u> </u>

FIG. 11D

	MCG1		MCC:		NCG3		MCG4	
DATA	Code Word	NCG	Code Word	NCG	Code Word MSB LSB	NCG	Code Word	NCG
-101	MSB LSB 000100010000010	2	0001000000010000	3	001001000100100	3	100000100000100	3
181 182	00010000100001	i	000010000001000	3	000100100100100	3	010001000010000	3
183	010001000100100	3	000001000000100	3	000010000010000	3	010000100001000	3
184	000001000000010	2 2	010000000100000	4	000001000001000	3	010000010000100	3
185	000000010000010		000000000100000	4	000000100000100	3	00100010001000	3
186	000100010010001	1	000000001000000	4	100100001000100	3	001000001000100	3
187	010010000010001	3	010010010000000 0100010010000000	4	100010001001000	3	100100000010000	3
188	010001000001000	2	010001001000000	4	100010000100100	3	100010000001000	3
189 190	·· 00010001001001	ī	010010001000000	4	001001001000100	3	100001000000100	3
191	010000100000100	3	010001000100000	4	001000100100100	13	010010000010000	3
192	000001000100010	2	001001001000000	4	000100000010000	3	010001000001000	3
193	000010000010010	2	001000100100000	4	000010000001000	3	001001000010000	3
194	000001000010001	1	010010000100000	4	100000001000000	4	001000100001000	3
195	000010000001001	1 3	001001000100000	4	100010001000000	4	001000010000100	3
196 197	001001001000100	3	000000100100000	4	1 100001000100000	4	010010001001000	3
198	001001000010000	3	001000100000000	4	100100001000000	4	010010000100100	3
199	000000100010010	2	000100010000000	4	100010000100000	4	001001001001000	3
200	0100001000000010	2	000010001000000	4	001001001000000	4	100100100010000	3
201	000000100001001	1	000001000100000	4	001000100100000	4	100100010001000	) š
202	010010001001001	1 2	010000100000000	4	001001000100000	4	100100001000100	3
203 204	000100000001000	3	000100001000000	4	000100100100000	4	100010010010000	] 3
205	001000010000010	2	000010000100000	4	000000100000000	4	100010001001000	1 3
206	0100000100000001	1	010000010000000	4	001000100000000	4	100010000100100	) 3
207	000100000010000	3	001000001000000	4	000100010000000	4	010010001000100	] 3
208	010010000100010	2	000100000100000	4	000001000100000	4	010001001001000	3
209	010001000010010	2	001000000100000	14	001000010000000	4	010001000100100	თოოოოოოოოოოოოოოოოოოოო <b>+</b> •
210 211	010001000001001	i	010000010010001	1	000100001000000	4	001001001000100	3
212	001000100001000	1 3	010000001001001	1	000010000100000	4	001000100100100	1 4
213	001001000001000	3	010000100010001	1	100000100000000	4	100000001000000	4
214	000100001000010	2	010000010001001	1 1	000100000100000	1 4	010000000100000	4
215	001001001001001	1 1	001000001001001	î	100000010000000	1 4	001000100100000	4
216 217	1:001000001000001 000010000001000	3	010001000010001	1	1 001000000100000	4	100100000100000	4
218	000010000000100	. 3	010000100001001	1	100001000010001	1	010010000100000	4
219	000010000100010	2	001000100010001	1	100000100001001	1 1	100000100000000	4
220	000100000100001	1	001000010001001	1	001000010010001	li	010000010000000	4
221	000010000010001		000100010010001	1 1	100010000010001	l i	001000001000000	4
222	000001000000100	3	0100000000001001	Î	100001000001001	1	100000010000000	4
223 224	000001000001001	Ĩ	010000000010001	1	001000100010001	1	010000001000000	4
225	010000010000000	4	001000000001001	1	001000010001001	1 1	001000000100000 100001000010001	li
226	001001000100010	2	001000000010001	1 1	000100010010001	$\begin{vmatrix} 1 \\ 1 \end{vmatrix}$	10000100001001	l î
227	001000100010010	2	000100000001001	$\begin{vmatrix} 1 \\ 1 \end{vmatrix}$	1000000000001001	Ιī	010000100010001	1
228	001001000010001	1	010010010001001	li	100000000010001	lī	010000010001001	1
229	001000100001001	3	000100000010001	lî	001000000001001	1	1 001000010010001	1 1
230 231	001000100000100	ı ă	000010000001001	١ī	100010010001001	1	001000001001001	1 1
232	000100010001000	13	010010000001001	1	100001001001001	1	100010000010001	1 1
233	001000001000000	4	001001000001001	1	001000000010001	1.1	01000100001001	li
234	000100001000100	3	000100100001001	1	000100000001001	1 1	010000100001001	î
<b>23</b> 5	000010010010000	3	000010010001001	1 1	10001001001001	li	001000100010001	1
236	000100000100000	4	000001001001001	li	000100000010001		001000010001001	. ( 1
237	000010001001000	3 4	0100010000010001	li	000010000001001	1 1	100100010001001	1
238 239	000100100100010	1 2	001001000010001	li	000010010001001	1	100010010010001	
240	000100010010010	2	001000100001001	1	000001001001001		100010001001001	-1-
240	1 moleculosicoro	شــــــــــــــــــــــــــــــــــــــ					<del></del> . <del></del> _	

## FIG. 11E

DATA	MCG1		MCG2		MCG3		MCG4	
SYMBOL	Code Word	NCG	Code Word MSB LSB	NCG	Code Word MSB LSB	NCC	Code Word MSB LSB	NCG
241 242 243 244 245	000100100010001 000100010001001 0000100010001000 001000000	1334	000100100010001 000100010011001 00001001	1 1 1 1	100100000010001 100010000001001 00100100	1	010010010001001 010001001001001 01001000000	1 1 1 1 1
246 247 248 249 250 : :	000010010001000	21323	000000011001001 00000010001001 00000010001001	1 1 1 1 1	000100010001001 000010010010001 000010001001	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	100010000001001 010010000010001 010001000010001 001001	1 1 1 1 1 1 1
252 253 254 255	000010001000100 000001001001000 00000001000000	3 3 1 3	010010001001001 001001001001001 000010000010001 000001000001001	1 1 1	100100001001001 001001001001001 000010000010001 000001000001001	1 1 1 1	100100001001001 010010010010001 010010001001	1 1 1 1 1

FIG. 12A

	DCG11	-	DCG12		DCG21		DCG22	
DATA SYMBOL	Code Word	NCG	Code Word	NCG	Code Word	NCG	Code Ford MSB LSB	NCG
3181000	MSB LSB	_	MSB 1.SB	-	01001001000000001	1	001001000000000001	1
000	00100100100000001	1	001001000000000001		01000100100000001	i	000100100000000000	1
001	00100010010000001	ì	000010010000000001	l i :	1010000100100000001	1	00001001000000001	1
002 003	001000001001000001	l î	000001001000000001	j	010000010010000001	1	00000100100000001	
003	00000100000000001	i	00000010010000001	1	01000000100100001	1	00000010010000001	1:
005	00100100010000001	1	00000001001000001	]	01001000100000001	] ]	00000000100100001	li
006	00100010001000001	1	00000000100100001	;	01000100010000001 01000010001000001		01000100000000001	li
007	00100001000100001	1	00100010000000001	1 1	01000001000100001	l i	00100010000000001	1
008	000100100100000001	1	0001000100000001		001001001000000001	l i	000100010000000001	1 1
009	00010001001000001	li	00000100010000001	i	100100010010000001	1	00001000100000001	1 !
010 T 011	00000010000000001	Ιi	00000010001000001		001000010010000001	1	00000100010000001	1 1
012	00100100001000001	l i	00000001000100001	1	00100000100100001	1	00000010001000001	li
013	100100010000100001	1	00100001000000001		000001000000000001 01001000010000001		010000100000000001	li
014	00010010001000001	1	00010000100000001		01000100001000001		00100001000000001	1
015	100010001000100001	1 1	00001000010000001	1 .	01000010000100001	] 1	1000100001000000001	13
016 017	00001001001000001	li	00000010000100001		00100100010000001	ן ו	00001000010000001	li
018	000000010000000001	ī	00100100100100001	1 1	00100010001000001		0000010000100001 00000010000100001	li
019	00000000001000001	]	00100000100000001		00100001000100001	$\mathbf{l}$	010010010010000001	1
020	00000000010000001	1 1	00010000010000001		000100010010000001	ī	01001000100100001	1
021	00100100000100001		00000100000100001		00010000100100001	1 1	010000010000000001	1 }
022 023	00001001000100001	lî	00100000000100001	1 1	000000100000000000	1	00100100100100001	1 1
024	00000100100100001		00100000001000001		01001000001000001		0010000010000001	l i
025	[00000000100000001	] ]	0001000000010000	1 2	01000100000100001	1 .	100001000001000001	.   1
026	000010000000000010	2 2	00100100000000010 000100100000000010		00100010000100001		100000100000100001	. ] 1
027 028	00100100100000010 00100010010000010	2	1000010010000000000	2	00010010001000001	1 1	01001001000100001	$\begin{pmatrix} 1 \\ 1 \end{pmatrix}$
029	00100001001000010	2	00000100100000010	) 2	0001000100010000		010001001001000001	1
030	100100000100100010	2	100000010010000010	) <u>2</u>	00001001001000001	i	00100000010000001	1
031	00100000010010010	2 2	00000001001000010		0000000100000000	il i	[0001000000100000]	
032 033	000001000000000010	2	00000000010010010		0000000001000000	1 1	00001000000100001	
034	00100010001000010	2 2 2	100100010000000000000000000000000000000	) 2	0100100000010000		01000000000100001	$\mathbf{i}$
035	10010001000100010	2	00010001000000010	2 2	0010010000010000	il î	1001000000000100001	1 1
036 -₹	00100000100010010	2 2	00001000100000010		0000100100010000	il ī	0100000001000000	1 1
037	00010010010000010		00000010001000010		10000010010010000	1 1	00100000001000001	
038 039	00010000100100010		00000001000100010	0 2	0000000010000000	1 1	00010000000100001 01001000000000001	2
040	100010000010010010	2	00000000100010010	2	0100100100000001	0 2	0010010000000000010	2
041	100000010000000010	2	0010000100000001	2 2	010001001000001		000100100000000010	
042	00000000001000010	2 2	0001000010000001		0100000100100001	012	1000010010000000010	3   2
043 044	00100100000010010		100000100001000010	0 2	[0100000010010001	0 2	00000100100000010	$\begin{array}{c c} 2 \\ 3 \\ 2 \end{array}$
045	00001001000010010	2	0000001000010001	0 2	0100000001001001	0 2	00000001001000010	
046	00000100100010010	2	0000000100001001	0 2	00001000000000001	0 2	0000000010010001	0   2
047	100000010010010010	1 2	001001001001000100010010010010010010010	$\begin{vmatrix} 2 \\ 0 \end{vmatrix} \begin{vmatrix} 2 \\ 2 \end{vmatrix}$	0100010001000001	ŏl ž	100000000010010010	0   2
048	00000000010000010	2	0010000010000001		[0100001000100001	0 2	01000100000000001	
049	00100100000100010	2	0001001001001001	0 2	in100000100010001	0  2	0010001000000001	0 2
-050 -051	00010010000100010	2	10001000001000001	0 2	0100000010001001	0 2	0001000100000001	
052	100010001000010010	) 2	0000100000100001	0 2	0010010010000001 0010001001000001	0  2 0  2	0000010001000001	0   2
053	00001001000100010	) 2	0000010700010001	0 2	Inninnoi100100001	01 Z	0000001000100001	0   2
054	00001000100010010	2 2	0000001300001001		ln010000010010001	0  Z	1 0000000100010001	0   2
055	0000010010010010010		0010000000010001	0 2	Inn10000001001001	01 2	0000000010001001	0 2
056 057	0000000100000010	) 2	100010000000001001	0 2	0000010000000001	0 2	0100001000000001	
058	100100100001000010	2	0010000000100001	0 2	0100100001000001	0 2	0001000010000001	0 2
059	100100010000100010	) 2	0001000000010001		I	0 2		0 2
060	00100001000010010	) 2	Imminimum	<u> </u>	10.100.00			

### FIG. 12B

DATA	DCG11		DCC12		DCG21		DCG22	
SYMBOL	Code Word	NCG	Code Word	NCG	Code Word MSB LSB	NCG	Code Word NSB LSB	NCG
061	00010010001000010	2	00100100100010010	-				↓
062	00010001000100010	2	0010010010010010	2	01000001000010010	2	00000100001000010	2
063	00010000100010010		00100000010000010	2	00100010001000010	2 2	00000010000100010	2
064	00001001001000010		00010000001000010	2	00100001000100010	2	01001001001000010	2 2
065	00001000100100010	2	00001000000100010	2	00100000100010010	2	01001000100100010	2
066	00001000010010010	2	00000100000010010	2	00010010010000010	2	01001000010010010	2
067	00100000000001000	3	001001000000000100	ã	00010001001000010	2	01000001000000010	2
068	000100000000000100	3	00010010000000100	3	00010000100100010	2	00100100100100010	2
069	00100000000010000	3	00001001000000100	3	00010000010010010	2	00100100010010010	5
070	00010000000001000	3	00000100100000100	3	000000100000000010	2	00100000100000010	2 2
071	000010000000000100	3	00000010010000100	3	01001000000010010	2	00010010010010010	2
072	00100100100000100	3	00000001001000100	3	001001000000010010	2	00010000010000010	2 2 2 2 2
073	00100010010000100	3	00000000100100100	3	00010010000010010	2 2	00001000001000010	2
074	00100001001000100	3	00100100000001000	3	00001001000010010	2	00000100000100010	2
075 076	1001000001001001001	3	00100010000000100	3	00000100100010010	2	00000010000010010	2
070	00010000000010000	3	000100100000001000	3	00000010010010010	2	01000000000010010	2
078	00000100000000100	3	00010001000000100   00001001000001000	3	000000000010000010	2	01000000000100010	2
079	00000000010010000	3	00001000100000100	3	01001000000100010	2 2	001000000000010010	2
080	00000000001001000	š	00000100100001000	3	010001000000010010 00100100000100010	2	00100000000100010	2
081	00000000100010000	ž	00000100010000100	3	00100010000010010	5	000100000000010010	5
082	00000000010001000	3	00000010010001000	ž	00010010000100010	2 2	01001001000010010	222222
083	00000000001000100	3	00000010001000100	3	00010001000010010	2	01000100100010010	2
084	00000001000010000	3	000000010010010001	3	00001001000100010	2	01000010010010010	2
085	00000000100001000	3	00000001000100100	3	00001000100010010	2	01000000010000010	2 2 2
086	00000000010000100	.3	00100100000010000	3	00000100100100010	2	00100000001000010	2
087	00100100010010000	3	00100010000001000	3	00000100010010010	2	000100000000100010	2
880	00100100001001000	3	00100001000000100	3	00000000100000010	2	00001000000010010	Z
08 <del>9</del> 090	00100100000100100	3	00010010000010000	3	01001000001000010	2	01001001000100010	2
091	000100100010010000	3	00010001000001000		01000100000100010	2	01001000100010010	2
092	00010010000100100	3	000010001000010000	3	01000010000010010  00100100001000010	2 2	01000100100100010   01000100010010010	2
093	00001001001001000	3	00001000100001000		00100010000100010	2	01000000100000010	2
094	00001001000100100		00001000010000100		00100001000010010	2	00100100100010010	2
095	00000100100100100	3	00001000100010000	3	00010010001000010		00100010010010010	2 2 2
095	00000010000010000		00001000010001000	3	00010001000100010	2	00100000010000010	
097	00000001000001000		00001000001000100	3	00010000100010010		00010000001000010	2
098	00000000100000100		00000100010010000		00001001001000010		00001000000100010	2 2
099 100	00100100100010000		00000100001001000		00001000100100010		000001000000010010	2
101	00100100010001000		00000100000100100		010000000000010000 001000000000001000		01001000000000100   001001000000000100	3 3
102	00100010010010000		00100000001001000		000100000000000100		000100100000000100	3
103	00100010001001000		00100000000100100		1001001000000100		000010010000000100	3
104	00100010000100100		00100000100010000		1000100100000100		00000100100000100	3
105	00010010010001000		00100000010001000	ž k	1000010010000100		00000010010000100	3
106	00010010001000100		00100000001000100		10000010010001001		00000001001000100	3
107	00010001001001000	3 0	00010000010010000		1000000100100100		00100100100100	3
108	00010001000100100		20010000001001000		0010000000000010000		010010000000001000	3
109	00001001001000100		000100000000100100		0001000000000001000		010001000000000100	3
	00001000100100100	3 10	00100001000010000		000010000000000100		00100100000001000	3
			00100000100001000		1001001000001000		00100010000000100	3
			00100000010000100		1001000100000100		00010010000001000	3
			00010000100010000		10001001000010001		00010001000000100	3
			0010000010001000		100010001000100	~ 1	0001001000001000	3
			000100001001000100		1000010010001000		000010010000100	3
			000010000010010000		10000100100100		00000100010000100	3
							00000010010001000	ž
119	00100001000100100						00000010001000100	ã
							0000001001001000	_

FIG. 12C

DATA	DCC11		DCG12		DCG21		DCG22	
SYMBOL	Code Word	NCG	Code Word	NCG	Code Word	NCG	Code Word NSB LSB	NCG
121	00010001001000100	3	00100001000001000	3	00100001001000100	3	00000001000100100	3
122	00010000100100100	3	00100000100000100		00100000100100100	Ĭš	01001000000010000	
123	00001000000010000	3	00010001000010000	3	00010000000010000	3	01000100000001000	
124	00000100000001000	3	00010000100001000	3	00001000000001000	3	010000100000000100	3
125	00000010000000100	3	00010000010000100	3	000001000000000100	3	00100100000010000	
126	00100000000100000	4	00000000001000000	4	00000000010010000	3	00100010000001000	
127	00000100100000000	4	00000000010000000	4	00000000100010000	3	00100001000000100	
128 129, .	00000010010000000	4	00100100100000000	4	00000000010001000	3	00010010000010000	
130	00000000100100000	4	00100010010000000	4	00000001000010000	3	00010001000001000 00010000100000100	3
131	00001000100000000	4	00100000100100000	4	00000000010001000	3	000100100001000	3
132	00000100010000000	4	00100100010000000	4	01001000010010000	3	00001000100001000	3
133	00000010001000000	4	00100010001000000	4	01001000001001000	3	00001000010000100	3
134	00000001000100000	4	00100001000100000	4	01001000000100100	3	00000100100010000	3
135	00010000100000000	4	000100100100000000	4	00100100010010000	3	00000100010001000	3
136	00001000010000000	4	00010001001000000	4	00100100001001000	3	00000100001000100	3
137 138	00000100001000000	4	00010000100100000	4	00100100000100100	3	00000010010010000	3
139	00100100100100000	4	00100010000100000	4	000100100100100000	3	00000100001001000  00000100000100100	3
140	00100000100000000	4	00010010001000000	4	00010010000100100	3	01000000010010000	3
141	00010000010000000	4	00010001000100000	4	00001001001001000	3	01000000001001000	3
142	00001000001000000	4	00001001001000000	4	00001001000100100	3	010000000000100100	3
143	00000100000100000	4	00001000100100000	4	00000100100100100	3	01000000100010000	3
144	001000000100000000	4	00100100000100000	4	00000010000010000	3	01000000010001000	3
)45 146	0001000000100000	4	00010010000100000	4	00000001000001000	3	010000000001000100	3
146	00100000001000000		00000100100100000	4	00000000100000100	3	00100000010010000	33
148	00010000000100000		00000000100000000	٦ I	01001000010001000	3	00100000000100100	3
149	01000000010010001		01001000010010001	i	01001000001000100	3	01000001000010000	3
150	01000000001001001		010010000001001001	1	01000100010010000	3	01000000100001000	3
151	01000000100010001		01001000100010001	3	01000100001001000	3	01000000010000100	3
152 153	01000000010001001		01001000010001001	1	01000100000100100	3	00100000100010000	3
154	01000001000010001 01000000100001001		01000100010010001   01000100001001001001	1	00100100100010000   00100100010001000	3	001000000100010001000 00100000001000100	3
155	01001001001001001		01000000000010001	i l	00100100001000100	3	00010000010010000	3
156 5.	01000010000010001		01001001000001001	ĩ l	00100010010010000		00010000001001000	3
157	01000001000001001		01000100100001001	1	00100010001001000		000100000000100100	3
158	01001000000001001		01000010010001001	1	00100010000100100		01001001001001000	3
159 - X 160	01001000000010001 010001000000001001		01000001001001001	1	00010010010001000		01001001000100100100	3
161	01000100000010001		01001001000010001 01001000100001001	1	00010010001000100		01000100100100100100 010000100000100001	3
162	01000010000001001		01000100100010001	i l	00010001000100100		01000001000001000	3
163	01001001000000001		010001000000000001	ī l	00001001001000100		01000000100000100	3
164	01000100100000001	1 0	010000100000000001	1	00001000100100100	3	00100001000010000	3
165	01000010010000001		1001001001000001	1	00000100000010000		00100000100001000	3
166	01000001001000001		1003000100100001	1	00000010000001000		00100000010000100	3
167 168	01000000100100001		1000001000000001	; [	00000001000000100		000100001000100000	3 3
169	01001000100000001		01001001000100001 01000100100100001	1	01001001000010000   0100100010000100001		000100000100010001 0001000000010001001	3
170	01000010001000001		1000000100000001	$\mathbf{i}$	01001000010000100		00001000010010000	3
171	01000001000100001	i	1000000000100001		01000100100010000	3	00001000001001000	3
172	01001000010000001	1 0	1000000001000001	i	01000100010001000	3 (	00001000000100100	3
	010010000000100001		1000000010000001	1	01000100001000100	3	01001001001000100	3
			10010000000000010		01000010010010000	3	01001000100100100	3
	01000100100000010		1000100000000010		01000010001001000		01000100000010000	3
	01000010010000010   01000001001000010		1000010000000010		01000010000100100		01000010000001000	3 3
			100100100100010		00100100100001000 00100100010000100		00100100100100100	3
			1001000010010010		00100010010001000	3	0100010000010000	3
			1000001000000010		00100010001000100	3 6	00100001000001000	3

# FIG. 12D

DATA	DCG11		DCG12		DCG21		DCG22	
SYMBOL	Code Word	NEC	Code Word	NCG	Code Word	NCG	Code Word	NCG
	MSB LSB	1	1120 1'20		MSB LSB		MSB LSD	
181 182	01000100010000010		0100000000010010	2	00100001001001000	3	00100000100000100	3
183	01000001000100010		0100000000100010		00100001000100100	3	00010001000010000	
184	01000000100010010		01001001000010010	2	00010001001000100	3	00010000010000100	3
185	01001000010000010		01000100100010010	2	00010000100100100	3	00001000100010000	3
186 187	01000100001000010		01000010010010010		00001000000010000	3	00001000010001000	3
188 -	01001000000010010   01001000000100010		01000000010000010		00000100000001000	3	00001000001000100	3
189 :-	01000100000010010		01001000100010010	2	01000000000100000	4	010000100100100000	3
190	01001000001000010		01000100100100010	2	01000000001000000	4	01000001001000000	4
191 192	01000100000100010	2	01000100010010010	2	00100000000100000	4	01000000100100000	4
193	01000010000010010 010000000000010000	3	01000000100000010	2	00000010010000000	4	01000100010000000 01000010001000000	4
194	01001001000000100	3	01001000000001000	3	00000000100100000	4	01000001000100000	4
195	01000100100000100		010001000000000100	3	00000100010000000	4	00100010010000000	4
196 197	01000010010000100		01001000000010000	3	00000010001000000	4	00100001001000000	4
198	01000001001000100	3	01000100000001000	3	000000010001000000	4	00100000100100000   01001000010000000	4
199	01001001000001000	3	01000000010010000	3	00000100001000000	4	010001000010000000	4
200	01001000100000100	3	01000000001001000	3	00000010000100000	4	01000010000100000	4
201 202	01001000010010000	3	01000000000100100	3	01001001001000000	4	00100100010000000	4
202	01001000001001000	3	01000000100010000	3	01001000100100000	4	00100010001000000	4
204	01001000100010000	š	01000000001000100	š	000100000100000000		000100100100000000	4
205	01001000010001000	3	01000001000010000	3	00001000001000000	4	00010001001000000	4
206 207	01001000001000100   01000100010010000	3	01000000100001000	3	00000100000100000	4	00010000100100000	4
208	0100010001001001000	3	01000000010000100     01001001001001000	3	01001001000100000   01000100100000		01001000001000000	4
209	01000100000100100	3	01001001000100100	ğΙ	00100000010000000		00100100001000000	4
210	01001001000010000	3	01000100100100100	3	00010000001000000		00100010000100000	4
211 212	01001000100001000   010010001000	3	01000010000010000   01000001000001000	3	00001000000100000		00010010001000000	4
213	01000100100010000	3	01000000100000100	3	01000000010000000 00100000001000000		00010001000100000	4
214	01000100010001000	3	01001001001000100	3	00010000000100000		00001000100100000	Å.
215 216	01000100001000100	3	01001000100100100	3	01000000010010001		01000000000010001	]
217	01000010010010000   01000010001001000	3	01000100000010000   01000010000	3	01000000001001001		00100000000001001 00100000000010001	1
218 🔩	D1000010000100100	3	01000001000000100	3	01000000010001001		000100000000001001	i
219	01000000000100000	4	01000100100000000		00100000010010001		01001001000001001	1
220 221	010000000010000001	4	01000010010000000		001000000001001001 01000001000010001		01000100100001001 01000010010001001	1
222	01001000100100000	4	01000000100100000		01000000100001001		01000001001001001	í
	01001001000100000	4	01001000100000000		00100000100010001	1 (	000100000000010001	1
	01000100100100000	4	01000100010000000		00100000010001001		000010000000001001	3
771	01000000100000000		01000010001000000 01000001000100000		00010000010010001 00010000001001001		01001001000010001	1
	00100000010010001		00000000010010001		01001001001001001		1000100100010001	i
228	00100000001001001	) l	000000000001001001		01000010000010001		1000100010001001	ī
	00100000100010001		00000000100010001		01000001000001001		1000010010010001	1
	001000000010001001 0001000001001001		000000000010001001	- 1	00100001000010001		01000010001001001 00100100100001001	1
	000100000001001001	• 1	00000001000010001		000100000100010001		00100010010001001	1
233	00100001000010001	1	00100100010010001		00010000010001001		0100001001001001	i
	00100000100001001		00100100001001001	1 (	00001000010010001		000010000000010001	ļ
	00010000100010001		00010010010010001 00010010001001001		000010000001001001		00000100000001001	1
237	00001000010010001		00001001001001001		00100100000001001	-	00000000100010001	1
238	000010000001001001	1  t	00000010000010001	1 (	00010010000001001	1 0	00000000010001001	i
	00100100000001001		00000001000001001		00001001000001001		00000001000010001	]
<u> 240 (</u>	00010010000001001		00100100100010001	3 (	00000100100001001	1 10	0000000100001001	1

### FIG. 12E

DATA	DCG11			DCG12		DCC	521		۵	CG22	
SYMBOL	Code Word	SB NCG	Code NSB	Word LSB	NCG	Code Wo MSB	LSB_	NCC	Code MSB	Word LSB	NCG
241	0000100100000010	)1 1	00100100	010001001	1	00000010010	0001001	1	010010000		1
242	0000010010000100	)1   1	00100010	010010001	1	00000001001	1001001	1	010010000		1
243	000100001000010	)] [	00100010	001001001	ì	01001000000	0010001	1	001001000		)
244	0000100010001000	)1   1	00010010	010001001	1	01000100000		1	001001000		)
245	0000100001000100	)1   1	00010001	100100100	1	00100100000		1	000100100		1
246	0000010001001001	)1   1	00000100	000010001	1	00100010000	0001001	1 1	000100100		l
247	0000010000100100	) 1   1	00000010	000001001	1	00010010000	0010001	1	000010010		1
248	00010001000000100	1 1	00100000	000001001	1	00010001000	0001001	1	000000100		1
249	00001001000001000	)1   1	00100000	000010001	1	00001001000	0010001	1	000000010		1
250	0000100010000100	1 1	00010000	000001001	1	00001000100	0001001	1	010010001		1
251	0000010010001000	1 1	00010000	000010001	1	00000100100	0010001	1	010010000		1
252	0000010001000100	1   1	00001000	000001001	1	00000100010	0001001	] ]	010001000		1
253	0000001001001000		00100100		1	00000010010		1	010001000		1
254	0000001000100100	1 1	00100010		1	00000010001		1	001001001		1
255	0010001000001000	1   1	00100001	1001001001	_l	01000100000	0010001	1	001001000	10001001	1

# FIG. 12F

DATA	DCG31		DCC32			DCG41		DCG42	
SYMBOL	Code Word	NCG	Code Word MSB LS	, NC	33	Code Word	NCC	Code Word	INC
000	10010010000000001	1	001001000000000		$\vdash$	MSB LSB	ļ	T NIZR TZR	
001	10001001000000001	l i	00010010000000000			100100100000000001	1	0010010000000000	
002	10000100100000001	l i	0000100100000000			10001001000000001	1	01000100000000000	
003	10000010010000001	l i	0000010010000000	i li		1000010010000001		0010001000000000	
004	10000001001000001	ì	0000001001000000	il i		10000001001000001	i	100001000000000001	
005	10000000100100001	1	00100010000000000	1 1		10000000100100001	i	00100001000000000	
006	10010001000000001	1	0001000100000000	ו ונ		10010001000000001	l i	10010010010000001	
007	10001000100000001	] ]	0000100010000000		- 1	10001000100000001	ì	10010001001000001	li
008 009	100000100010000001	1	0000010001000000			10000100010000001	1	10010000100100001	l ī
010	1000001000100001	1 1	1000001000100000			10000010001000001	1	100000100000000001	
Ŏij	10010000100000001	il	10000100000000000	-		10000001000100001	1	01001001001000001	1
012	10001000010000001		0001000010000000			01001001000000001 01000100100000001	1 1	01001000100100001	1 1
013	10000100001000001	ī	0000100001000000			01000010010000001	l i	01000001000000001 00100100100100001	1
014	10000010000100001	1	0000010000100000			01000001001000001	l i	00100000100000001	1
015	[00100100100000001]		0000001000010000	1 1		1000000100100001	î	1.0010010001000001	li
016	00100010010000001	1	1001001001000000	1 1	- [:	10010000100000001	1	10010001000100001	li
017 018	00100001001000001	1	1001000100100000	1 1		10001000010000001	1	100010010010000001	۱ī
019	00100000100100001 000001000000000001	1	1001000010010000	1 1		10000100001000001	1	10001000100100001	1
020	10010000010000001	i b	10000010000000000 001001001001001000	$\begin{array}{c c} 1 & 1 \\ 1 & 1 \end{array}$		10000010000100001	1	10000001000000001	1
021	10001000001000001	î li	0010000010000000	il i		)1001000100000001  )1000100010000001	1	01001001000100001	1 1
022	10000100000100001	i l	0001000001000000	il i		1000010001000001	1	010001001001000001 010000001000000001	1 :
023	00100100010000001		0000100000100000			1000001000100001	î	001000000100000001	1 1
024	00100010001000001		0000010000010000	1 1		0100100100000001	i	10000000000100001	l i
025 026	00100001000100001		1001001000100(100		10	0100010010000001	1	10000000001000001	Ī
027	00010010010000001 00010001001000001		1001000100010000	-		0100001001000001	1	01000000000100001	1
028	00010000100100001		1000100100100100000 10001000100100000			0100000100100001	1	10000000010000001	1
029	00000010000000001		000000100000000			0010000010000001 0001000001000001	1 1	01000000001000001	1
030	10010000001000001		010000001000000			0000100000100001	i l	001000000000100001 10010010000100001	1
031	10001000000100001		0010000001000001			1001000010000001	i l	10001001000100001	i
032	00100100001000001		0001000000100001			1000100001000001	ī	10000100100100001	î
033 034	00100010000100001	; l;	0000000000100001	1		0010000000100001	1	10000000100000001	Ĩ
035	00010010001000001	1	00000000001000001	1 1		1001000000100001	1	010000000010000001	1
			0000000010000001 0100000000100001			0100100000100001	1	001000000001000001	Ţ
			0010010000100001			0010010000000010 00010010000000010	2	010010000000000010   001001000000000010	2 2
038			0001001000100001			0000100100000010		100010000000000010	ž
		1   1	0000100100100001	l i		0000010010000010	2	010001000000000010	2
			0000000100000001	1	10	0000001001000010		001000100000000010	2
		i lõ	0100000001000001	1		0000000100100010	2	100001000000000010	2
		2 0	0010000000100001 0100100000000000010	1 2		0000000010010010	2	010000100000000010	2
			001001000000000000000000000000000000000	2		00100010000000010 00010001000000010		0010000100000000010	2
			00010010000000000	2				100100100100000010   100100010010000010	2 2
046		2 0	0000100100000001010	2		000010001000010		10010000100100010	2
	10000001001000010	2 00	00000100100000010	2			_ 1	10010000010010010	2
	10000000100100010	2   10	00010000000000010	2				100000100000000010	2
	100000000010010010		1000100000000010	2				01001001001000010	2
	100100010000000010  2 10001000100000010  2		0100010000000010			000300300000	- 1	01001000100100010	2
	0001000100000010  2  0000100010000010  2		0001000100000010 1000100010000010					01001000010010010	2
	0000010001000010 2		0000100010000010					010000010000000010	2
	0000001000100010		00001000000000010					00100100100100010   00100100010010010	2
555 1	0000000100010010 2		1000010000000010					0100000100000010	2
	200100000000010 2	00	010000100000010				-	00000000000100010	2
	0010000100000010 2		001000010000010	2	10	000100001000010		1000000000010010	2
	0001000010000010 2		000100001000010				2   1	.00000000001000010	2
	0000100001000010 2 0000010000100010 2		000010000100010					1000000000100010	2
1 <u>1 00 </u>	0000010000100010 2	110	010010010000010	2	010	0010001000000010  :	2   0	0100000000010010	2

FIG. 12G

DATA		DCG31					CG32			DCG	41			DCG42	
SYMBOL	Code IISB	Word L.	SB	4CC	MSB	Code	Word LS8	NCG	Code MSB	· Wo	rd LSB	NCG	Code NSB	Word LSB	NCG
061	10000001			2	1001	00010	01000010	2	01000100	010	000010	2	100100100	000010010	2
062	00100100			2			00100010	2	01000010	2001	000010	2	100010010	000010010	2
063	00100010			2	1001	00000	10010010	2	01000001	1000	100010	2	100001001	100010010	2
	00100001			2	1000	00100	00000010	2	101000000	100	010010	2	100000100	10010010	1 2
	00100000			2	0010	01001	00100010	2	00100100	100	000010	2	100000000	10000010	2 2
	001000000			2	0010	01000	10010010	2	00100010			2	010000000	01000010	2
	000001000			2	0010	00001	000000010	2	00100001	1001	000010	2	001000000	000100010	2222222222222222222
	100100100			2	0001	20100	10010010	2	00100000	100	100010	2	100100100	000100010	2
	100100000	01000001	10	2	00010	00000	10000010	2	00100000	010	010010	2	100100010	000010010	2
070	100010000	00100001		2	0000	10000	01000010	2	10010010	010	010010	2	100010010	01000100	2
071	100001000	00010001	0	2	00000	01000	00100010	2	10010000	010	000010	2	100010001	00010010	2
	100000100			2	00000	00100	00010010	2	10001000	0001	000010	2	100001001		2
	001001000			2			00100010	2	10000100	0000	100010	2	100001000	10010010	2
	001000100	00100001		2	10000	00000	01000010	2	10000010	0000	010010	2	1.000000001	.00000010	2
	001000010			2	00100	00000	00010010	2	01001000	010	000010	2	010010010	000010010	2
	001000001			2	10010	20100	00010010	2	01000100	001	000010	2	010001001		2
	000100100			2	1000	10010	00010010		10010000			2	010000100		2
	000100001			2	10000	10010	00010010		01001000	0000	010010	2	010000000		2
	000010010		0	2			10010010		00100100			2	001000000		2
	000010001			2			10000010		10010000			2	100100100		2
	000010000						00100010	2	10001000			2	100100010		2
	100100000						00010010	2	01001000			2	100100001		2
	001001000						0100010	2	01000100			2	100010010		2
	000100100		0 3				0010010	2	00100100			2	100010001		2
	000010010						0100010		00100010			2	100010000		2
	000001001						0010010		10010000			2	100000010		2
	000000100						0100010	2	10001000			2	010010010		2
	100100000						0010010	2 2	10000100			2	010010001		Z
	100010000						0000010	2	01001000			2	010001001		2
	001001000						1000010	2	01000100			5	010001000		2
	001000100						0100010	Ž	01000010			2	010000001		2
	000100100						0010010	2	00100100			2	001001001		2
	000100010						1000010	2	00100010			2	001000100		2
	000010010						0100010		00100001			2 3	100100000		2
	000010001						0010010	51	10010010			3	010010000		3
	000001001 000001000						1000010   0100010		10000100			3	001001000		222233333
	100100000						0010010		10000100	7000	20100	3	100100000		ă
	1000100000						0000010		10000010			3	100010000		3
	1000010000						0010010		100000000			3	010010000		3
	001001000								010000000			3	010001000		3
	00100 <b>010</b> 00						0010010		001000000			3	001001000		3
	0010000100						1000010		100100100			3	001000100		3
103 11	<b>MINNOIN</b>	っついていいてん	, L	. [	JULIU	لالحماد		4	TOUTOUTO	بالحالات		٦	1001000100		จั

## FIG. 12J

DATA	DCG31		DCG32		DCG41			CG42	
SYMBOL	Code Word NSB LSB	NCG	Code Word	NCG	Code Word MSB LSB	NCG	Code		NCG
241 242 243 244 245 246 247 248 249 250	10000001000001001 00100000100010001 001000000	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	NSB	1 1 1 1 1 1 1 1 1	MSB	1 1 1 1 1 1 1	MSB 010001001 010000100 010000010 100100001 100010001 100010000 100001000 010010	1000 100 1 0100 100 1 000 100 1 0000 100 1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
251 252 253 254 255	00100000100001001 00010000100010001 000100001001	1 1 1 1 1	00100000000010001 100000100100001001 1000001001	1 1 1 1 1 1 1 1 1	01000100100001000 01000100010000100 0100001001	3333	010010010 010010001 01001000 010001000	00010000 00001000 10000100 00010000	33333

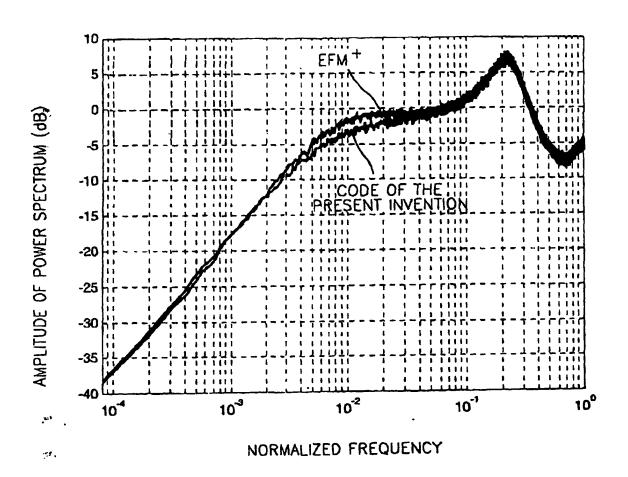
FIG. 13A

		ACG1		ACG2		VCC3		ACG4	
DATA SYMBOL	Code		NCG	Code Word	NCG	Code Word	NCG	Code Word MSB LSB	NCG
000 001 002 003 004 005 006 007 008 009 010 013 014 015 016 017 018 020 021 022 023 024 025 026 027 028 029 030 031 032 033 034 035 037 039 030 031 032 033 034 035 036 037 038 039 039 039 039 039 039 039 039 039 039	00000100 0000100 0100100 0100100 0100100	00000000 00000000 00000000 0000000 00000	44441312	00001000000000 00001000000000 001001000000	44442	100100010000001 10010010000001 100100100	111122222222233333333333333333333333333	000001000000001 00001000000001 000010001000001 000010001000001 000010001000001 000010001000001 000010001000001 000010001000001 000010001000001 000010001000001 000010001000001 000010001000001 000010001000001 00001000100001 00001000100001 00001000100001 00001000100001 00001000100010001 00001000100010001 00001000100010001 00001000100010001 0000100010001001	111111111111111111111111111111111111111

FIG. 13B

DATA	<u> </u>		ACG 1				ACG2			ACG3		İ	ACG4	
YMBOL	IISB	Code	Word	LSB	NCG	Code MSB	Word LSB	NCG	Code NSB	Word LSB	NCG	Code NSB	Kord	NC
061 062	l											00010000		3
063						ı		1 }			1 1	00010000	0010000	3334
064								1			1 1	00010010	0100100	1 3
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FIG. 14





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(11) **EP 1 251 641 A3** 

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### **EUROPEAN PATENT APPLICATION**

(88) Date of publication A3: 16.03.2005 Bulletin 2005/11

(51) Int Ci.7: H03M 5/14

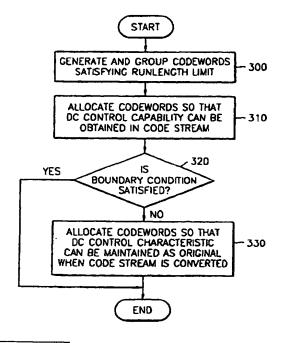
- (43) Date of publication A2: 23.10.2002 Bulletin 2002/43
- (21) Application number: 02252790.7
- (22) Date of filing: 19.04.2002
- (84) Designated Contracting States:
  AT BE CH CY DE DK ES FI FR GB GR IE IT LI LU
  MC NL PT SE TR
  Designated Extension States:
  AL LT LV MK RO SI
- (30) Priority: 20.04.2001 KR 2001021360
- (60) Divisional application: 04010409.3 / 1 450 491 04010410.1 / 1 450 492
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### (54) Code generation and allocation method

(57)A method for generating and allocating codewords is provided. The method includes allocating one of two selectable codewords b1 and b2 as codeword b when a preceding codeword a and a following codeword b form a code stream X, in which codewords b1 and b2 have opposite INVs which are parameters indicating whether the number of '1s' contained in a codeword is an odd number or an even number and when the code stream of a and b1 is X1, and the code stream of a and b2 is X2, allocating codewords such that the INVs of X1 and X2 are maintained to be opposite when a or b1(b2) should be replaced by another codewords in compliance with a predetermined boundary condition given between codewords. According to the method, by using a short codeword having less bits as a main conversion codeword, high efficiency is achieved in recording density. Also, when codewords which do not satisfy the run length conditions are replaced by other codewords, the codewords are allocated so that the DC suppression capability of the code stream can be maintained, and therefore higher DC suppression capability of the code stream is provided.

FIG. 3



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### EUROPEAN SEARCH REPORT

Application Number EP 02 25 2790

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